

Space-Efficient IBE without Pairings

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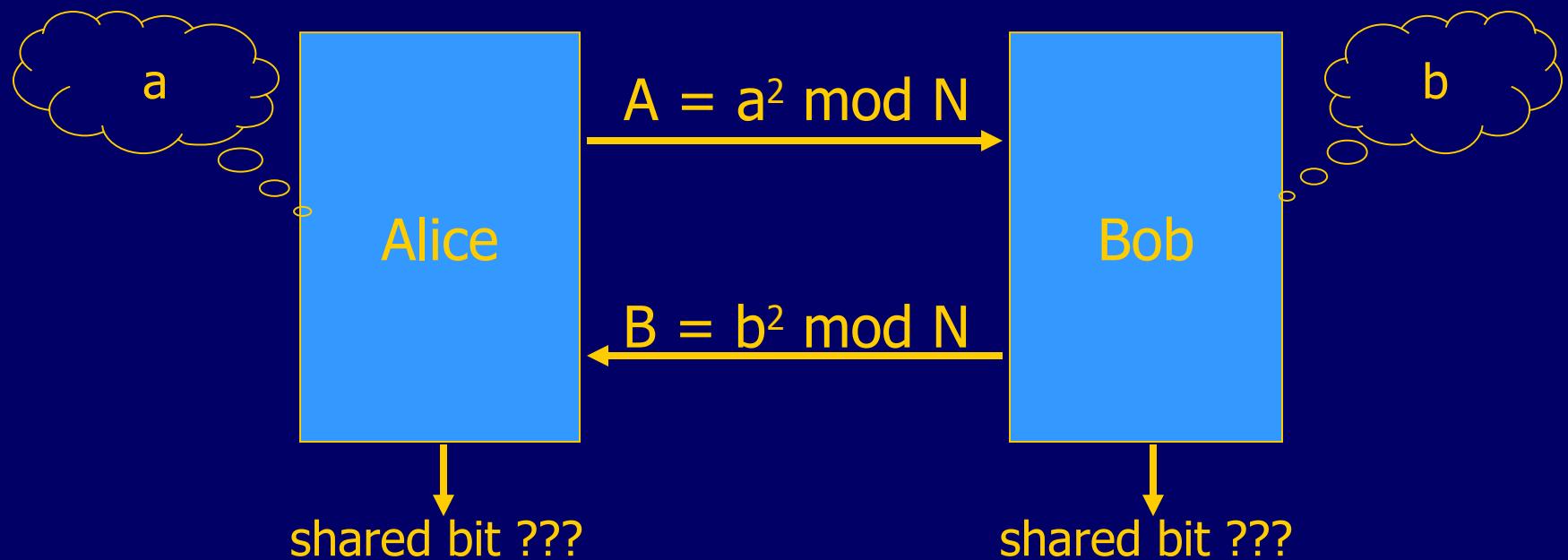
Stanford University

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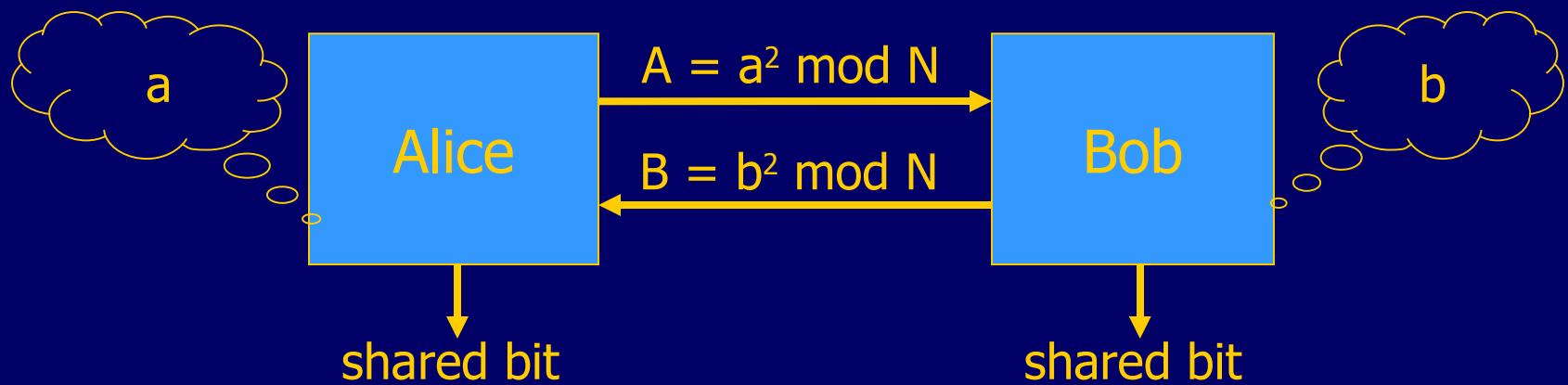
IBE without Pairings

- Most work on IBE uses pairings (a.k.a. bilinear maps):
 - Boneh-Franklin (2001), BB04a, BB04b, Wat05, Gen06, etc.
- An IBE Scheme by Clifford Cocks (2001)
 - No pairings!
 - Based on quadratic residuosity mod N (w/ R.O.)
 - ... but CT is 2×1024 times longer than PT
- New system: $|CT| = |PT| + \log_2 N + 1$

Prelude to our scheme: 1-bit Key Agr.



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$$A \cdot x^2 + B \cdot y^2 = 1 \text{ mod } N$$

$$\text{Shared bit} = \text{Jac}(ax+1, N) = \text{Jac}(2by+2, N)$$

$$A \cdot x^2 + B \cdot y^2 = 1 \pmod{N} \quad (*)$$

- Ong-Schnorr-Shamir (1984) signature scheme:

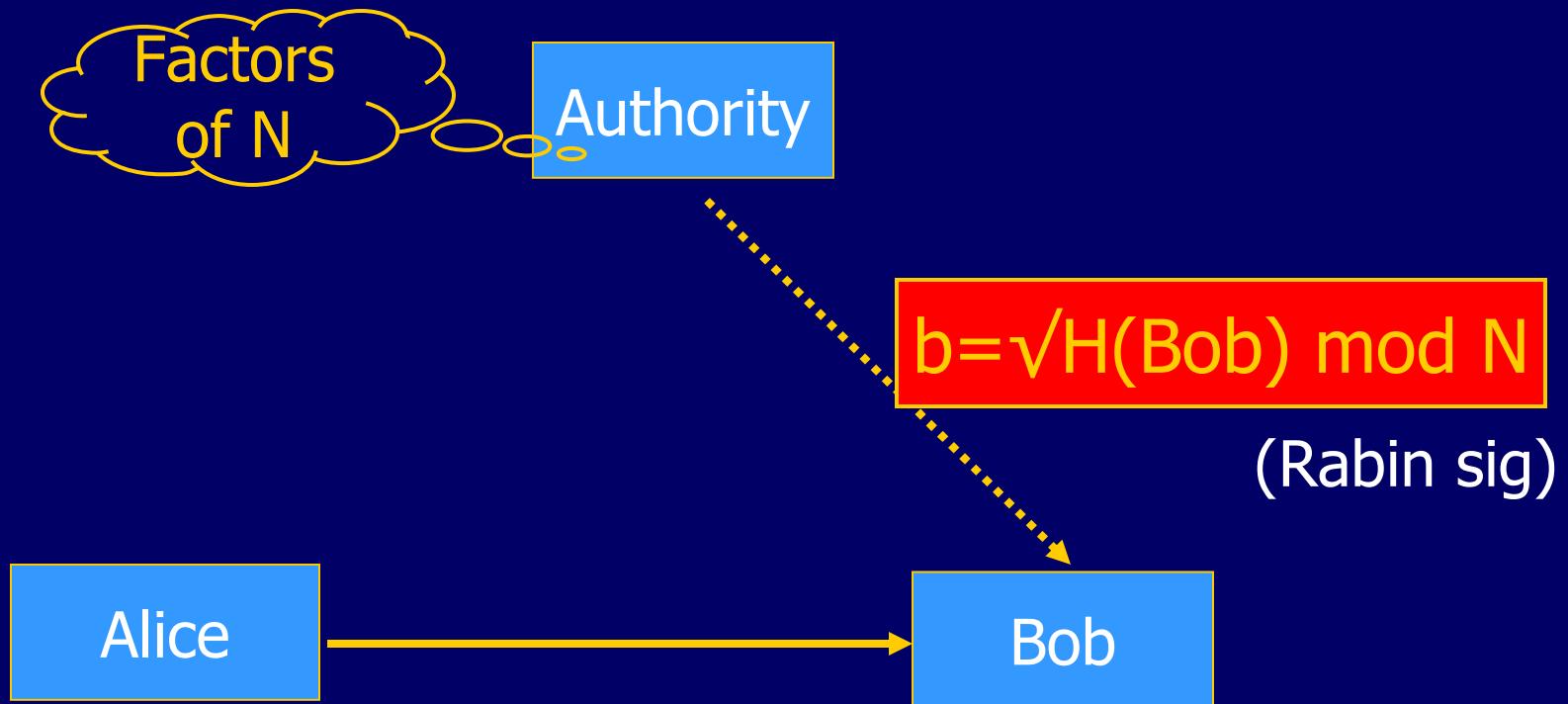
For $\text{PK}=(N, A)$,

$\text{sig}(m) = (x, y)$ s.t.

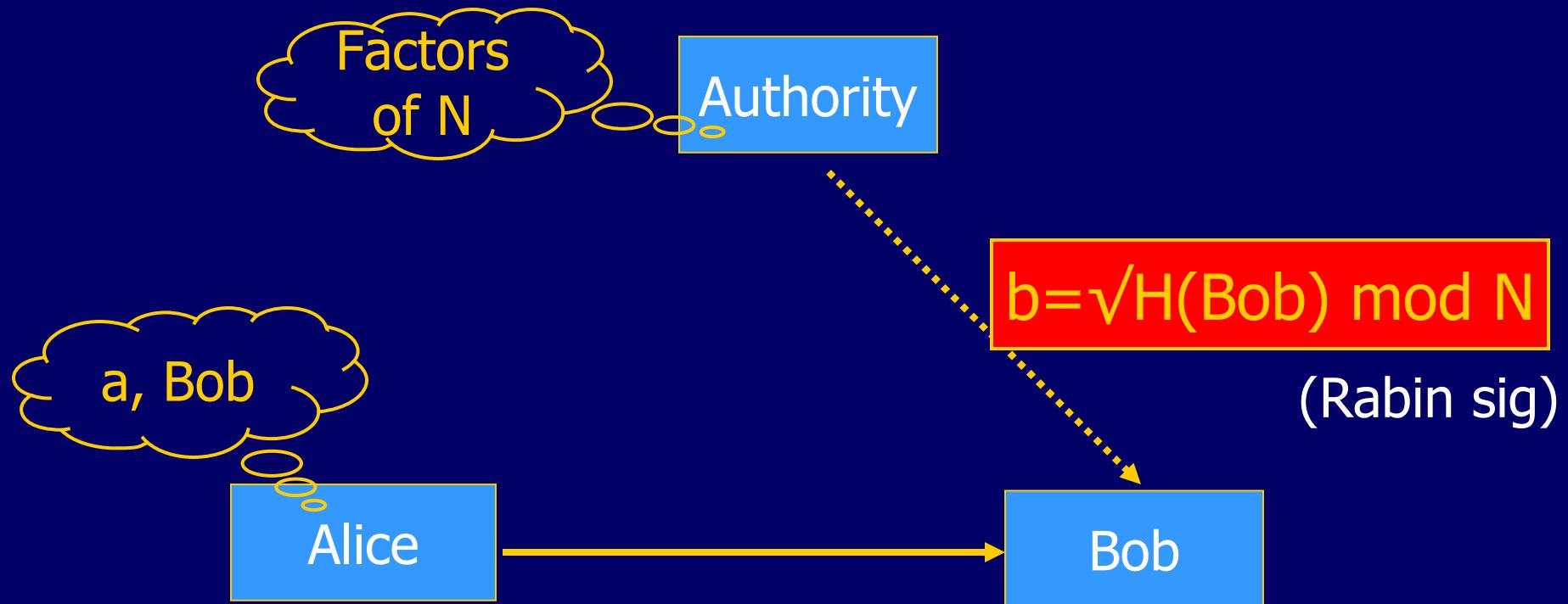
$$A \cdot x^2 + m \cdot y^2 = 1 \pmod{N}$$

- ... but efficient algorithm to solve $(*)$
Pollard-Schnorr algorithm
- So, our 1-bit key agreement scheme works

The 1-bit IBE system



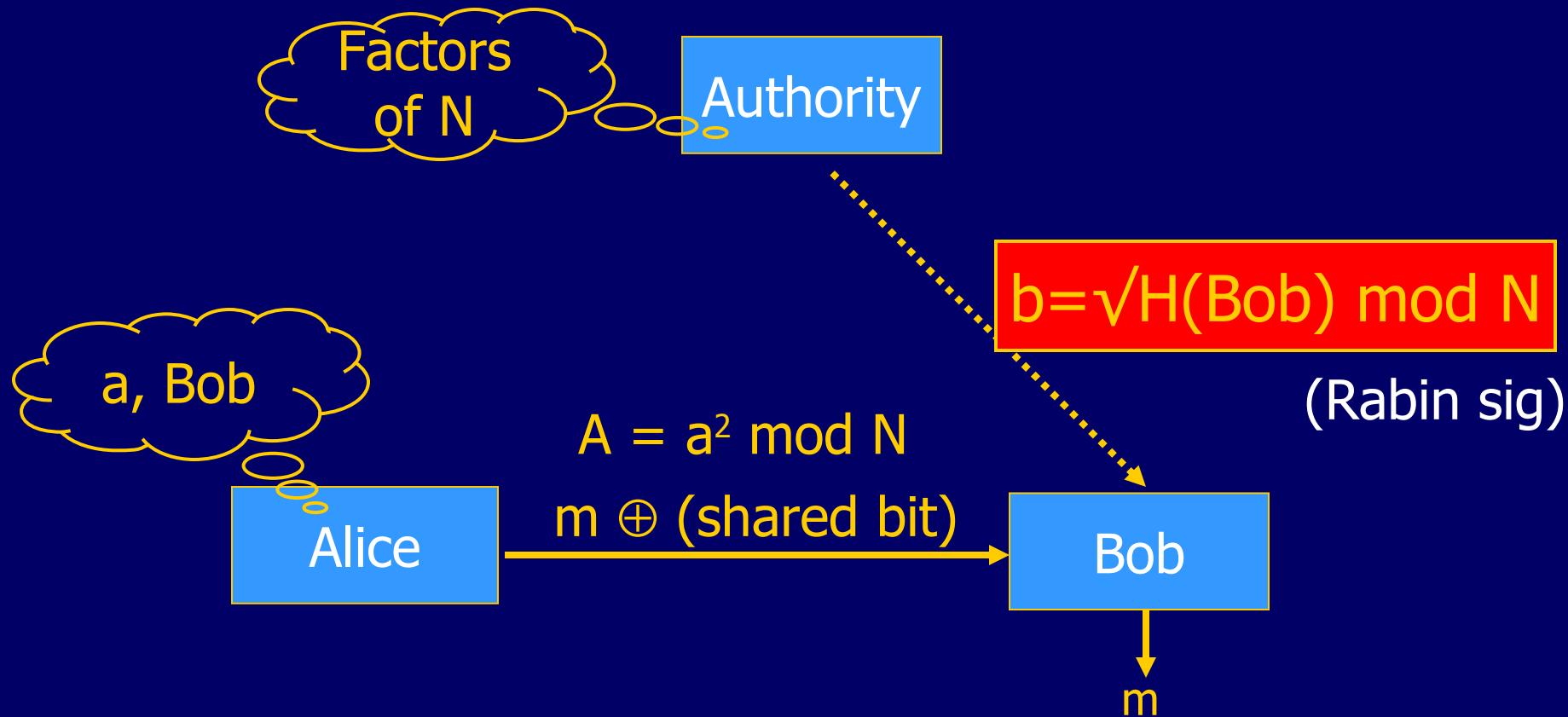
The 1-bit IBE system



The 1-bit IBE system

(x,y) s.t.

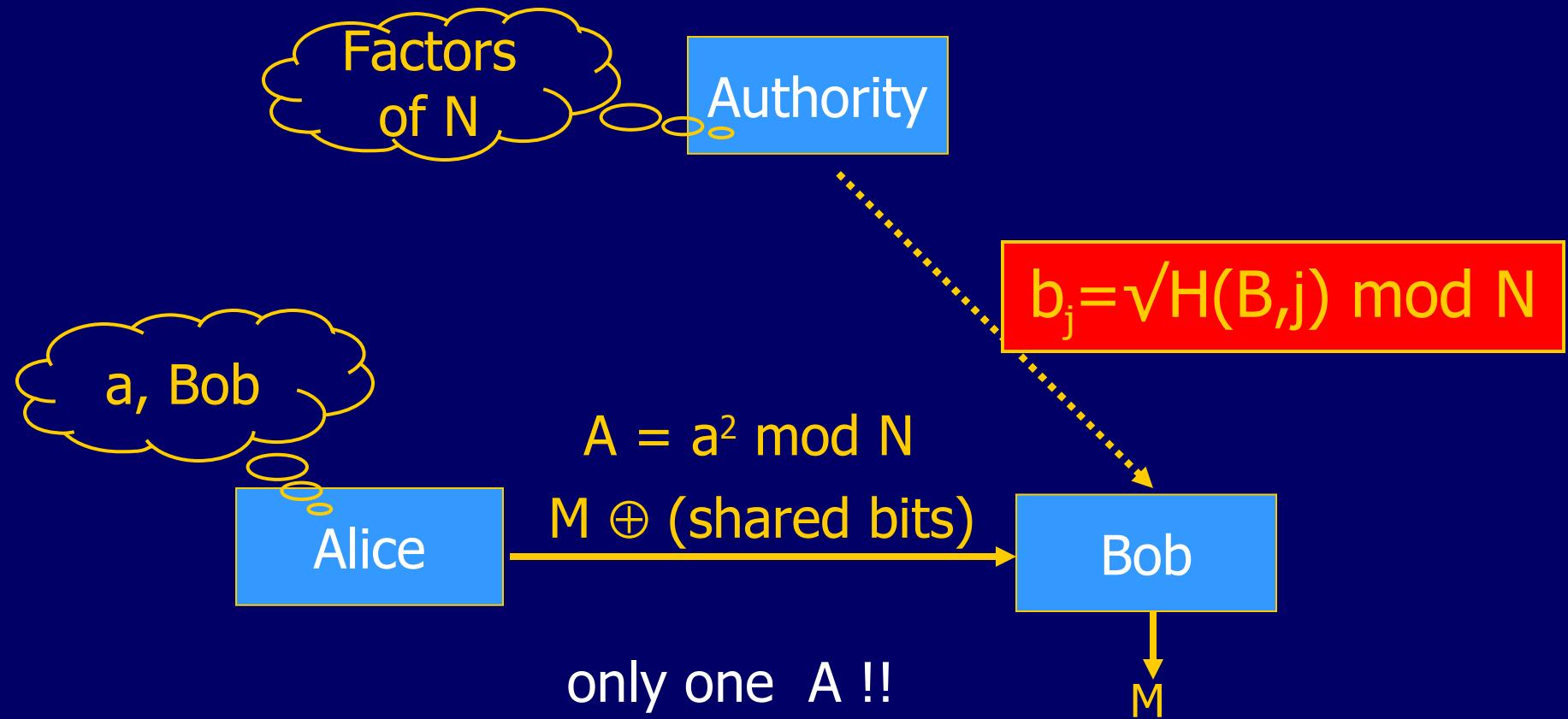
$$A \cdot x^2 + H(\text{Bob}) \cdot y^2 = 1 \pmod{N}$$



Multi-bit IBE system

(x,y) s.t.

$$A \cdot x^2 + H(B,j) \cdot y^2 = 1 \pmod{N}$$



Security, Extensions, Open Problems

- Security:
 - ANON-IND-ID-CCA based on QR (w/ R.O.)
 - R.O. needed for Rabin sigs security
- Extensions:
 - New hash proof system based on QR
 - New standard model PKE system based on QR
- Open Problems:
 - Solve $Ax^2 + By^2 = 1 \text{ mod } N$ faster!
 - Remove random oracles